

Optimized Analytical Processing New Features with 11g R2

Hüsnü Şensoy

Global Maksimum Data & Information Technologies

Founder, VLDB Expert

husnu.sensoy@globalmaksimum.com

Agenda

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 - Understanding Optimized Analytical Processing Capabilities
- New Capabilities by 11g Release 2
 - Multi-Predicate Partition Pruning
 - Intelligent Multi-Branch Execution
 - NULL Aware ANTI JOIN
 - Hash-Based DISTINCT Aggregation
- Conclusion

Who am I ?

- Data & Information expert on VLDB environments
 - DWH
 - Data Mining
 - Inference Systems
 - Data Archiving Solutions
 - Niche Storage Technologies
 - Recovery Strategies & Solutions
 - HA Systems
- Oracle ACED on BI field
 - Only one in Turkey
 - Still the youngest one all over the community.
- DBA of the Year 2009
 - 7th and still the youngest all over the world.
 - Only one in Turkey
- Blogger @ <http://husnusensoy.wordpress.com>
- Member of Oracle CAB for 12g DWH development
- Worldwide presenter of Oracle conferences and user group events



Optimized Analytical Processing New Features with 11g R2

Introduction

Optimized Analytical Processing Capabilities

- *Optimized Analytical Processing Capabilities* are those features implemented by Oracle on CBO, SQL execution, and expression manager that transparently improve SQL performance for your data crunching processes.
- The keyword is transparency. In many circumstances, you don't need to change any configuration to enable those capabilities.
- Oracle keeps saying «*SQL is X times faster in this release*» mainly due to those features.
- It is usually very hard to hear about them until the product is mature or some of them cause problems in your production.

Why should I care about them ?

- In 10g one of the most important headaches for large DWH customers was related to new hash group by optimizations. Many customers have disabled them with the guidance of support (`_gby hash aggregation enabled`). So being familiar with new SQL engine will let you a better understanding of product and give you the chance to take remedial actions.
- Most people are annoyed with SQL plan changes with each release. They usually choose to freeze them using various techniques. Understanding those new features will let you to understand the reasons behind some plan changes in new release.
- Just to appreciate the effort made by those developers optimizing our lives.

Optimized Analytical Processing New Features with 11g R2

Multi-Predicate Partition Pruning

Partition Pruning

- In one of recent surveys, Oracle partitioning seems to be the *Top 1* feature used by large DWH sites.
- Range partitioning not only helps ILM in DWH but also improves query performance by partition pruning most of the time.
- Until 11gR2 Oracle is *biased* on using *static partition* pruning rather than *dynamic* one.
- *Multi-predicate pruning* is the idea of utilizing each and every possible pruning opportunity to reduce the amount of data to be read from disk or buffer cache.

Partitioning Scheme for *SH.SALES*

SH.SALES

1995-1996

One partition per
year

1997

One partition for
each half of a year

1998-2003

One partition for
each quarter of a year

A Simple Query on *SH.SALES*

```
select /*+ FULL (s) FULL (t) */ count (*)
  from sh.sales s, sh.times t
 where s.time_id = t.time_id
 and t.fiscal_month_name in ('February')
 and s.time_id between
      to_date ('01-JAN-1998', 'DD-MON-YYYY')
 and
      to_date ('01-JAN-2001', 'DD-MON-YYYY');
```

Partition Pruning Opportunities on SH.SALES

```
select /*+ FULL (s) FULL (t) */ count (*)  
  from sh.sales s, sh.times t  
 where s.time_id = t.time_id  
 and t.fiscal_month_name in ('February')  
 and s.time_id between  
      to_date('01-JAN-1998', 'DD-MON-YYYY')  
 and  
      to_date('01-JAN-2001', 'DD-MON-YYYY');
```

Pruning Idea	Description	Partitions to be scanned on SH.SALES
No Pruning		Scan all 28 partitions.
Static Pruning	Use predicate on <i>time_id</i> column of <i>SH.SALES</i>	Scan only 13 partitions.
Dynamic Pruning	Build a filter list for month <i>February</i> on <i>SH.TIMES</i> table then access to <i>SH.SALES</i> table.	Scan only 5 partitions
Static + Dynamic Pruning (Multi-predicate Pruning)	Use static & dynamic pruning together.	Scan only 3 partitions

In 10.2.0.4

Execution Plan

Plan hash value: 68236240

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		1	24	217 (11)	00:00:03		
1	SORT AGGREGATE		1	24				
*2	HASH JOIN		41164	964K	217 (11)	00:00:03		
*3	TABLE ACCESS FULL	TIMES	84	1344	9 (0)	00:00:01		
4	PARTITION RANGE ITERATOR		684K	5344K	202 (9)	00:00:03	5	17
*5	TABLE ACCESS FULL	SALES	684K	5344K	202 (9)	00:00:03	5	17

A Simple Query on *SH.SALES* with Tracing Add-ons

```
set autot trace exp stat
alter session set tracefile_identifier = 'multiPredicatePruning';
alter session set events '10128 trace name context forever, level 2';

select /*+ FULL (s) FULL (t) */ count(*)
  from sh.sales s, sh.times t
 where s.time_id = t.time_id
 and t.fiscal_month_name in ('February')
 and s.time_id between
      to_date('01-JAN-1998', 'DD-MON-YYYY')
 and
      to_date('01-JAN-2001', 'DD-MON-YYYY');

alter session set sql_trace = false;
```

Execution Plan in 11.2.0.1

Execution Plan

Plan hash value: 3278936322

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		1	24	322 (8)	00:00:05		
1	SORT AGGREGATE		1	24				
*2	HASH JOIN		43252	1013K	322 (8)	00:00:05		
3	PART JOIN FILTER CREATE	:BF0000	91	1456	13 (0)	00:00:01		
*4	TABLE ACCESS FULL	TIMES	91	1456	13 (0)	00:00:01		
5	PARTITION RANGE AND		690K	5393K	303 (7)	00:00:05	KEY (AP)	KEY (AP)
*6	TABLE ACCESS FULL	SALES	690K	5393K	303 (7)	00:00:05	KEY (AP)	KEY (AP)

10128 Trace Content in 11.2.0.1

...

Partition Iterator Information:

partition level = PARTITION

call time = RUN

order = ASCENDING

Partition iterator for level 1:

iterator = ARRAY [count= 3, max = 28] = **4 8 12**

...

Remarks

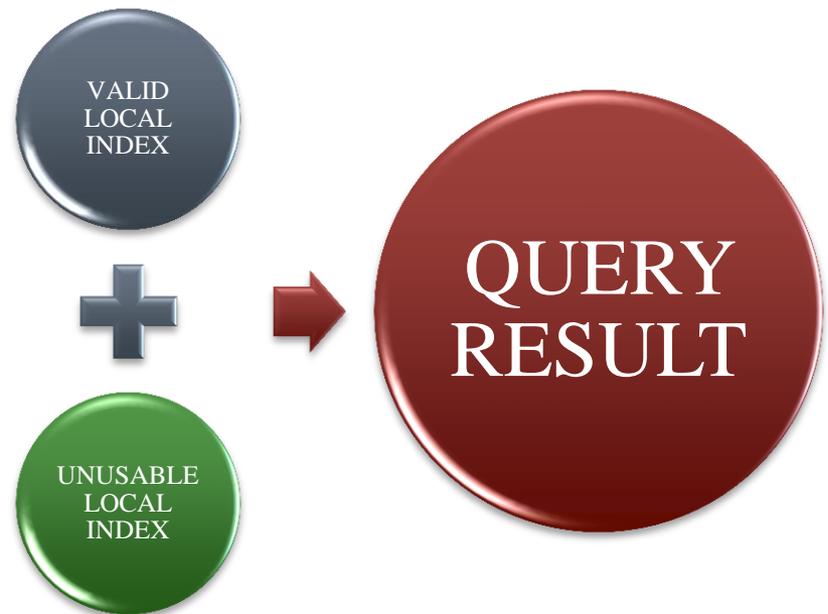
- Partitioning is and will be *Number One* trick of very large data management and processing.
- *Multi-predicate Partition Pruning* boosts Oracle's pruning opportunities for cases where several predicates can result in pruning.

Optimized Analytical Processing New Features with 11g R2

Intelligent Multi-Branch Execution

Horizontal Partial Indexing

- As you may all know, Oracle allows *UNUSABLE* index partitions starting from early releases of partitioning technology.
- Many data warehouses wish to disable some old index partitions to reveal the burden of maintaining them during ELT.
- Intelligent Multi-Branch Execution is a query rewrite technique to split a single SQL statement based on a partitioned table having LOCAL indices.



Another Simple Query on *SH.SALES*

```
select channels.channel_desc,  
       sum(sales.amount_sold) as total_amount  
from sh.sales, sh.products, sh.channels  
where channels.channel_id = sales.channel_id  
and products.prod_id = sales.prod_id  
and channels.channel_class = 'Direct'  
and products.prod_category = 'Photo'  
group by channels.channel_desc  
order by 2 desc;
```

Execution Plan

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		2	66	581 (2)	00:00:07		
1	SORT ORDER BY		2	66	581 (2)	00:00:07		
2	HASH GROUP BY		2	66	581 (2)	00:00:07		
* 3	HASH JOIN		29505	950K	575 (1)	00:00:07		
* 4	TABLE ACCESS FULL	CHANNELS	2	42	3 (0)	00:00:01		
5	PARTITION RANGE ALL		70812	829K	571 (1)	00:00:07	1	28
6	TABLE ACCESS BY LOCAL INDEX ROWID	SALES	70812	829K	571 (1)	00:00:07	1	28
7	BITMAP CONVERSION TO ROWIDS							
8	BITMAP AND							
9	BITMAP MERGE							
10	BITMAP KEY ITERATION							
11	BUFFER SORT							
12	TABLE ACCESS BY INDEX ROWID	PRODUCTS	14	294	2 (0)	00:00:01		
* 13	INDEX RANGE SCAN	PRODUCTS_PROD_CAT_IX	14		1 (0)	00:00:01		
* 14	BITMAP INDEX RANGE SCAN	SALES_PROD_BIX					1	28
15	BITMAP MERGE							
16	BITMAP KEY ITERATION							
17	BUFFER SORT							
* 18	TABLE ACCESS FULL	CHANNELS	2	42	3 (0)	00:00:01		
* 19	BITMAP INDEX RANGE SCAN	SALES_CHANNEL_BIX					1	28

Disable A Few LOCAL Index Partitions

```
ALTER INDEX SH. SALES_CHANNEL_BIX MODIFY PARTITION SALES_1995 UNUSABLE;  
ALTER INDEX SH. SALES_CHANNEL_BIX MODIFY PARTITION SALES_1996 UNUSABLE;  
ALTER INDEX SH. SALES_CHANNEL_BIX MODIFY PARTITION SALES_H1_1997 UNUSABLE;  
ALTER INDEX SH. SALES_CHANNEL_BIX MODIFY PARTITION SALES_H2_1997 UNUSABLE;  
ALTER INDEX SH. SALES_CHANNEL_BIX MODIFY PARTITION SALES_Q1_1998 UNUSABLE;  
ALTER INDEX SH. SALES_CHANNEL_BIX MODIFY PARTITION SALES_Q2_1998 UNUSABLE;  
ALTER INDEX SH. SALES_CHANNEL_BIX MODIFY PARTITION SALES_Q3_1998 UNUSABLE;  
ALTER INDEX SH. SALES_CHANNEL_BIX MODIFY PARTITION SALES_Q4_1998 UNUSABLE;
```

Execution Plan in 10.2.0.4

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		2	108	278 (15)	00:00:03		
1	SORT ORDER BY		2	108	278 (15)	00:00:03		
2	HASH GROUP BY		2	108	278 (15)	00:00:03		
* 3	HASH JOIN		75870	4000K	267 (11)	00:00:03		
4	MERGE JOIN CARTESIAN		24	1008	5 (0)	00:00:01		
* 5	TABLE ACCESS FULL	CHANNELS	2	42	3 (0)	00:00:01		
6	BUFFER SORT		14	294	2 (0)	00:00:01		
7	TABLE ACCESS BY INDEX ROWID	PRODUCTS	14	294	1 (0)	00:00:01		
* 8	INDEX RANGE SCAN	PRODUCTS_PROD_CAT_IX	14		0 (0)	00:00:01		
9	PARTITION RANGE ALL		910K	10M	254 (9)	00:00:03	1	28
10	TABLE ACCESS FULL	SALES	910K	10M	254 (9)	00:00:03	1	28

Execution Plan in 11.2.0.1

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		63850	1558K	386 (6)	00:00:06		
1	SORT ORDER BY		63850	1558K	386 (6)	00:00:06		
2	HASH GROUP BY		63850	1558K	386 (6)	00:00:06		
3	VIEW	VW_TE_12	63850	1558K	376 (4)	00:00:06		
4	UNION-ALL							
* 5	HASH JOIN		44707	1790K	273 (2)	00:00:04		
* 6	TABLE ACCESS FULL	CHANNELS	2	42	3 (0)	00:00:01		
7	PARTITION RANGE ITERATOR		53648	1047K	269 (2)	00:00:04	9	28
8	TABLE ACCESS BY LOCAL INDEX ROWID	SALES	53648	1047K	269 (2)	00:00:04	9	28
9	BITMAP CONVERSION TO ROWIDS							
10	BITMAP AND							
11	BITMAP MERGE							
12	BITMAP KEY ITERATION							
13	BUFFER SORT							
14	TABLE ACCESS BY INDEX ROWID	PRODUCTS	14	294	2 (0)	00:00:01		
* 15	INDEX RANGE SCAN	PRODUCTS_PROD_CAT_IX	14		1 (0)	00:00:01		
* 16	BITMAP INDEX RANGE SCAN	SALES_PROD_BIX					9	28
17	BITMAP MERGE							
18	BITMAP KEY ITERATION							
19	BUFFER SORT							
* 20	TABLE ACCESS FULL	CHANNELS	2	22	3 (0)	00:00:01		
* 21	BITMAP INDEX RANGE SCAN	SALES_CHANNEL_BIX					9	28
* 22	HASH JOIN		19143	1159K	98 (10)	00:00:02		
23	MERGE JOIN CARTESIAN		24	1008	5 (0)	00:00:01		
* 24	TABLE ACCESS FULL	CHANNELS	2	42	3 (0)	00:00:01		
25	BUFFER SORT		14	294	2 (0)	00:00:01		
26	TABLE ACCESS BY INDEX ROWID	PRODUCTS	14	294	1 (0)	00:00:01		
* 27	INDEX RANGE SCAN	PRODUCTS_PROD_CAT_IX	14		0 (0)	00:00:01		
28	PARTITION RANGE ITERATOR		229K	4486K	90 (7)	00:00:02	1	8
29	TABLE ACCESS FULL	SALES	229K	4486K	90 (7)	00:00:02	1	8

Remarks

- *Intelligent Multi-Branch Execution* is an invaluable new optimization for sites using LOCAL indexes.
- Keep in mind in order to use this optimization SKIP_UNUSABLE_INDEXES parameter should set to be TRUE.
- This option can be disabled by setting _OPTIMIZER_TABLE_EXPANSION parameter to FALSE.

Optimized Analytical Processing New Features with 11g R2

NULL Aware ANTI-JOIN

ANTI JOIN

- Oracle can use ANTI JOIN execution plan (with Nested Loop, Hash, or Merge join options) in case that a SQL statement contains NOT IN or NOT EXISTS clauses (or something rewritten to this).
- Hash Anti-Join is known to be an optimal execution plan for many large data warehouse queries containing above clauses.
- One major problem about classical anti-join is that due to some design errors like constraint ignorance, Oracle will reject using anti-join (not to generate erroneous result sets) and put a FILTER step instead (Refer [one of my earlier blog posts](#)).
- FILTER is usually CPU consuming and never-ending step for the join of large datasets.

Yet Another Simple Query on *SH.SALES*

```
select count(*) from sh.sales
      where time_id not in (select time_id
                             from sh.times);
```

Anti Join

Execution Plan

Plan hash value: 397380204

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		1	16	41 (22)	00:00:01		
1	SORT AGGREGATE		1	16				
* 2	HASH JOIN RIGHT ANTI		9188	143K	41 (22)	00:00:01		
3	INDEX FAST FULL SCAN	TIMES_PK	1826	14608	3 (0)	00:00:01		
4	PARTITION RANGE ALL		918K	7178K	29 (0)	00:00:01	1	28
5	BITMAP CONVERSION TO ROWIDS		918K	7178K	29 (0)	00:00:01		
6	BITMAP INDEX FAST FULL SCAN	SALES_TIME_BIX					1	28

Release NOT NULL Constraint on *SH.SALES*

```
alter table SH.SALES modify TIME_ID NULL;
```

Execution Plan in 10.2.0.4

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		1	8	5670 (2)	00:01:01		
1	SORT AGGREGATE		1	8				
* 2	FILTER							
3	PARTITION RANGE ALL		910K	7112K	29 (0)	00:00:01	1	28
4	BITMAP CONVERSION TO ROWIDS		910K	7112K	29 (0)	00:00:01		
5	BITMAP INDEX FAST FULL SCAN	SALES_TIME_BIX					1	28
* 6	INDEX FULL SCAN	TIMES_PK	1	8	4 (0)	00:00:01		

Predicate Information (identified by operation id):

```
2 - filter( NOT EXISTS (SELECT /*+ */ 0 FROM "SH"."TIMES" "TIMES" WHERE LNNVL("TIME_ID"<>:B1)))
6 - filter(LNNVL("TIME_ID"<>:B1))
```

Execution Plan in 11.1.0.6+

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		1	16	41 (22)	00:00:01		
1	SORT AGGREGATE		1	16				
* 2	HASH JOIN RIGHT ANTI SNA		9188	143K	41 (22)	00:00:01		
3	INDEX FAST FULL SCAN	TIMES_PK	1826	14608	3 (0)	00:00:01		
4	PARTITION RANGE ALL		918K	7178K	29 (0)	00:00:01	1	28
5	BITMAP CONVERSION TO ROWIDS		918K	7178K	29 (0)	00:00:01		
6	BITMAP INDEX FAST FULL SCAN	SALES_TIME_BIX					1	28

Create SH.SALES_NEW & SH.TIMES_NEW

```
create table sh.sales_new nologging
tablespace users as
select      s.*,
           to_char(time_id,'DD') day_id,
           to_char(time_id,'MM') month_id,
           to_char(time_id,'YYYY') year_id
from sh.sales s;
```

```
create table sh.times_new
nologging tablespace users as
select      t.*,
           to_char(time_id,'DD') day_id,
           to_char(time_id,'MM') month_id,
           to_char(time_id,'YYYY') year_id
from sh.times t;
```

```
alter table sh.sales_new modify day_id not null;
alter table sh.sales_new modify month_id not null;
alter table sh.sales_new modify year_id not null;
```

```
alter table sh.times_new modify day_id not null;
alter table sh.times_new modify month_id not null;
alter table sh.times_new modify year_id not null;
```

How about this Query ?

```
select count(*)  
from sh.sales_new  
where (day_id,  
       month_id,  
       year_id) not in (select day_id,  
                           month_id,  
                           year_id  
                        from sh.times_new);
```

Simple Execution Plan

Execution Plan

Plan hash value: 3458658284

```
-----  
| Id | Operation | Name | Rows | Bytes | Cost (%CPU) | Time |  
-----  
| 0 | SELECT STATEMENT | | 1 | 22 | 1527 (1) | 00:00:19 |  
| 1 | SORT AGGREGATE | | 1 | 22 | | |  
|* 2 | HASH JOIN RIGHT ANTI | | 9188 | 197K | 1527 (1) | 00:00:19 |  
| 3 | TABLE ACCESS FULL | TIMES_NEW | 1826 | 20086 | 17 (0) | 00:00:01 |  
| 4 | TABLE ACCESS FULL | SALES_NEW | 918K | 9870K | 1507 (1) | 00:00:19 |  
-----
```

Remove Only One NULL Constraint on *SH.TIMES*

```
alter table sh.times_new modify year_id null;
```

Execution Plan (NA)

Pre 11g Release 2

By 11g Release 2

Execution Plan

Plan hash value: 1287449578

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time
0	SELECT STATEMENT		1	10	1502K (1)	05:00:27
1	SORT AGGREGATE		1	10		
* 2	FILTER					
3	TABLE ACCESS FULL	SALES_NEW	756K	7384K	1505 (1)	00:00:19
* 4	TABLE ACCESS FULL	TIMES_NEW	1566	15660	2 (0)	00:00:01

Execution Plan

Plan hash value: 1313396486

Id	Operation	Name	Rows	Bytes	TempSpc	Cost (%CPU)	Time
0	SELECT STATEMENT		1	22		5545 (1)	00:01:07
1	SORT AGGREGATE		1	22			
2	MERGE JOIN ANTI NA		9188	197K		5545 (1)	00:01:07
3	SORT JOIN		918K	9870K	42M	5527 (1)	00:01:07
4	TABLE ACCESS FULL	SALES_NEW	918K	9870K		1507 (1)	00:00:19
* 5	SORT UNIQUE		1826	20086		18 (6)	00:00:01
6	TABLE ACCESS FULL	TIMES_NEW	1826	20086		17 (0)	00:00:01

Remove Only One NULL Constraint on *SH.SALES*

```
alter table sh.times_new modify year_id not null;  
alter table sh.sales_new modify year_id null;
```

Execution Plan (SNA)

Pre 11g Release 2

By 11g Release 2

Execution Plan							Execution Plan							
Plan hash value: 1287449578							Plan hash value: 1679400675							
Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Id	Operation	Name	Rows	Bytes	TempSpc	Cost (%CPU)	Time
0	SELECT STATEMENT		1	10	1502K (1)	05:00:27	0	SELECT STATEMENT		1	22		5545 (1)	00:01:07
1	SORT AGGREGATE		1	10			1	SORT AGGREGATE		1	22			
* 2	FILTER						2	MERGE JOIN ANTI SNA		9188	197K		5545 (1)	00:01:07
3	TABLE ACCESS FULL	SALES_NEW	756K	7384K	1505 (1)	00:00:19	3	SORT JOIN		918K	9870K	42M	5527 (1)	00:01:07
* 4	TABLE ACCESS FULL	TIMES_NEW	1566	15660	2 (0)	00:00:01	4	TABLE ACCESS FULL	SALES_NEW	918K	9870K		1507 (1)	00:00:19
							* 5	SORT UNIQUE		1826	20086		18 (6)	00:00:01
							6	TABLE ACCESS FULL	TIMES_NEW	1826	20086		17 (0)	00:00:01

Remove Only One NULL Constraint on *SH.SALES* & *SH.TIMES*

```
alter table sh.times_new modify year_id null;  
alter table sh.sales_new modify year_id null;
```

Execution Plan (NA)

Pre 11g Release 2

By 11g Release 2

Execution Plan

Plan hash value: 1287449578

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time
0	SELECT STATEMENT		1	10	1502K (1)	05:00:27
1	SORT AGGREGATE		1	10		
* 2	FILTER					
3	TABLE ACCESS FULL	SALES_NEW	756K	7384K	1505 (1)	00:00:19
* 4	TABLE ACCESS FULL	TIMES_NEW	1566	15660	2 (0)	00:00:01

Execution Plan

Plan hash value: 1313396486

Id	Operation	Name	Rows	Bytes	TempSpc	Cost (%CPU)	Time
0	SELECT STATEMENT		1	22		5545 (1)	00:01:07
1	SORT AGGREGATE		1	22			
2	MERGE JOIN ANTI NA		9188	197K		5545 (1)	00:01:07
3	SORT JOIN		918K	9870K	42M	5527 (1)	00:01:07
4	TABLE ACCESS FULL	SALES_NEW	918K	9870K		1507 (1)	00:00:19
* 5	SORT UNIQUE		1826	20086		18 (6)	00:00:01
6	TABLE ACCESS FULL	TIMES_NEW	1826	20086		17 (0)	00:00:01

Remarks

- *NULL Aware ANTI JOIN* is a great enhancement for constraint ignorant databases.
- SNA is first introduced in 11g Release 1, but multi column support is now available by 11g Release 2
- SNA is not a way to cheat SQL design practices.
- This option can be disabled by setting `_optimizer_null_aware_antijoin` parameter to `FALSE`
- To learn more about *NULL Aware ANTI JOIN*, refer to [great post by Greg Rahn](#).

Optimized Analytical Processing New Features with 11g R2

Hash-Based DISTINCT Aggregation

HASH GROUP BY

- After 10g Oracle starts to use HASH GROUP BY instead of SORT GROUP BY more extensively as it is appropriate.
- This is fundamentally related with hashing has a lower time complexity ($O(n)$) than sorting ($O(n \log n)$).
- DISTINCT clause inhibits Oracle from using HASH GROUP BY and force it to utilize SORT GROUP BY instead.
- And some unlucky Telco customers heavily utilizes DISTINCT COUNT clause in their queries (number of distinct subscribers).

Yet Another Simple Query on *SH.SALES*

```
select sum(QUANTITY_SOLD) total_sold ,  
       count(distinct channel_id) ndiff_channel  
from sh.sales  
group by prod_id;
```

Pre 11.2.0.1 Execution Plan

Execution Plan

Plan hash value: 4109827725

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		72	720	515 (7)	00:00:07		
1	SORT GROUP BY		72	720	515 (7)	00:00:07		
2	PARTITION RANGE ALL		918K	8973K	488 (2)	00:00:06	1	28
3	TABLE ACCESS FULL	SALES	918K	8973K	488 (2)	00:00:06	1	28

Execution Plan by 11.2.0.1

Execution Plan

Plan hash value: 913412106

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		72	2160	515 (7)	00:00:07		
1	HASH GROUP BY		72	2160	515 (7)	00:00:07		
2	VIEW	VW_DAG_0	204	6120	515 (7)	00:00:07		
3	HASH GROUP BY		204	2040	515 (7)	00:00:07		
4	PARTITION RANGE ALL		918K	8973K	488 (2)	00:00:06	1	28
5	TABLE ACCESS FULL	SALES	918K	8973K	488 (2)	00:00:06	1	28

Be Careful !!!

```
select sum(QUANTITY_SOLD) total_sold ,  
       count(distinct channel_id) ndiff_channel,  
       count(distinct time_id) ndiff_time  
from sh.sales  
group by prod_id;
```

Even in 11.2.0.1

Execution Plan

Plan hash value: 4109827725

Id	Operation	Name	Rows	Bytes	Cost (%CPU)	Time	Pstart	Pstop
0	SELECT STATEMENT		72	1296	515 (7)	00:00:07		
1	SORT GROUP BY		72	1296	515 (7)	00:00:07		
2	PARTITION RANGE ALL		918K	15M	488 (2)	00:00:06	1	28
3	TABLE ACCESS FULL	SALES	918K	15M	488 (2)	00:00:06	1	28

Remarks

- I believe, this feature have no customer coverage as much as others but if you are one of those *distinct counters*, you will definitely benefit from it.
- Actually the part I have introduced is a part of all hash group by optimizations introduced with 11g Release 2. For appropriate use of all optimizations you might need to fix *Bug 9148171* in 11.2.0.1.
- More than one distinct count do not work.
- This option can be disabled by setting `_optimizer_distinct_agg_transform` parameter to FALSE.

Conclusion

- There are many more optimized analytical processing capabilities introduced in Oracle 11g Release 2.
- Those are all about fine tuning the existing features instead of introducing new fancy ones.
- And to be honest that's what large customers want.



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